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A Multi Level Zone Based Congestion Control Mechanism for Multicast Routing in MANETS

A.Mallaiah[1] , P.Imrankhan[2], Dr. G.N.Swamy[3]

[1]*Associate Professor, Dept.of ECE, Gudlavalleru Engineering College, Gudlavalleru, A.P, INDIA*

[2]*M.Tech (DECS) student, Dept.of ECE, Gudlavalleru Engineering College, Gudlavalleru, A.P, INDIA*

[3]*Dr.G.N.Swamy, Professor, VR Siddardha Engineering College, Vijayawada, A.P, INDIA*

Abstract--- **In this paper A MLZBCC mechanism overcrowding discovery method is projected. The prepared model intends to allocate an energy efficient mechanism to compute the amount of blockage at casualty node with maximal accuracy. This congestion detection equipment is incorporated with a Two Step Multi degree area established Congestion Control Routing Protocol. The proposed model entails controlling of blockage in two steps with successful energy capable blocking detection and optimum utilization of resources. Packet loss in network routing is mainly due to link failure and blockage. A lot of the present congestion control options would not possess the capability to distinguish between packet loss because of link collapse and packet loss as a result of congestion. As a consequence these remedy objective towards activity against packet drop because of link failure which is an effort and may merchandise in loss of resources. The other perimeter in most of the accessible option is the use of power and resources to discover blockage state, measure of congestion and alarm the origin node about blocking in routing route. Here within this paper we suggest multiple apparatus model of congestion recognition an control apparatus that include energy-efficient congestion detection, Zone level Congestion Evaluation Algorithm [ZCEA] and Zone level Egress Regularization Algorithm [ZERA].This document is reinforced by the investigational and simulation results demonstrate that better cost minimization, energy-efficiency in congestion control and congestion detection is achievable by the proposed method.**

***Keywords*—Ad-hoc networks, multiple-mechanism design, optimization, random access, wireless networks**.

I. INTRODUCTION

The normal TCP jamming control largely designed for web is really not an pertinence for MANETs because MANETs are proven to influence protocol and protocol heaps of control systems and also the MANETs are environmentally irreconcilable with standard TCP .The package saving delays and losses in MANETs are mostly as a result of their node mobility combined with fundamentally unforeseen medium which really is actually a direct consequence of the usual wireless multiple hop channel can't be construe as congestion losses .The primary identity of a wireless multiple hop route is that within interfering range of one node just a single data is transmitted. In Magnet's sites in a complete place are overloaded due to common where as net blockage is single router. A note valuable point is the fact that in a MANET the nodes are not overcrowded.

It's tough to locate the source of blockage in a multi-hop network because one consumer has the ability to generate a congestion causing relatively lower bandwidth of mobile ad-hoc networks. The wifi networks are somewhat more prone to congestion problems compared with the conventional wire line system. Consequently a well-balanced congestion control system is to be employed compulsorily for the stability and outstanding performance of the wireless network. The non-homogeneous feature of the program protocols in the multihop wifi networks, one and unified solution for the congestion related problems can't be found. Instead an appropriate congestion control based upon the attributes and functions of the related network can be designed. As a result, these proposition majorly kind a subset of remedies for the determined problems rather than a complete, promptly employed protocol. They present like a parent for program-tailored protocol stacks. Exceptionally, some of the protocol properties serve broad variety of applications.

The recent years have witness a much more concentrate in the congestion control methods directed to the modeling, analysis, algorithm development of closed loop control techniques (e.g. TCP) making them considerate for adaption to the cellular hoc systems. Under the provision of constraints of routing route and bandwidth calculations possessing the capacity to unify and support operation have been evolved. Another important constraint to be painstaking in a wireless hoc community is due to the MLZBCC [Multi Stage Zone Based Congestion Control) mechanism. Bulks of wireless MLZBCC possess a moment constraint letting an individual user to obtain a physical channel at a specified time.

As to the parts within the document are organized to offer the following details. The section2 explores the most reported works within the region of text, section3 provides a depth discussion of the projected protocol and section 4 depends on the simulator and their results to become consummate by conclusion and recommendations.

II. RELATED WORK

Hierarchical centric playing control solution is found in [1].Analytics based alternative for congestion aware routing was offered in [4], [2] introduced metrics to evaluate data rate, MLZBCC foil and buffer delay, which helps to distinguish and offer the obstructing contention region in network. HongqiangZhai [3] in the offing a remedy by arguing that congestion and severe medium discourse is interrelated. Yung Yi [4] suggested a jump level obstruction control model. Tom Goff, Nael [5] mentioned a set of calculations that initiates different path use once the feature of the path being used becomes suspect. Xuyang [6] present a multiple-mechanism hop-by-hop congestion control scheme planned to enhance TCP performance in multi-hop wireless networks. The impact is presented by dzmitry et al [7] preventing on transport device that decreases the operation. Duc [8] proposed that existing designs for routing aren't flexible to blockage.

The existent models aim at determine congestion losses in routing route. The package loss creates a link failure. Making attempts to control the packet losses that cause link failure are in successful. Yet another exclusive approach is regularizing the outlet at all nodes participating in routing. In most cases of control the congestion at hop level [15]. Henceforth egress regularization at every node of the system necessitates operation of pricey riches. Here in this paper we argue it is a very important to recognize the reason for packet loss. Hence we may prevent the blockage control method via egress regularization below the position of link failure. As well as we continue the spat that jump stage blocking control alone isn't plenty if the hop levels are unable to stabilize themselves. The weight to handle the blocking by using the same resources can be done as in springtime level egress regularization models. Within this case in our prior work we proposed a Two Step numerous Mechanisms preventing Control Routing Protocol [18]. Here we suggest a new energy-efficient multiple device based preventing control routing method that contains Congestion recognition and congestion control models. Multiple mechanism Design with Control plane functionality in Multi hop Wireless Networks for efficient blocking Control.

A. Energy Resourceful Congestion Detection Mechanism

The objective of the proposed congestion detection device is always to capture amount of congestion at exchange hop degree node with maximal accuracy. In model, the detection mechanism is decoupled from various other activities of the MLZBCC mechanism including link consistency analysis and stream size analysis. The acknowledgement model expanded to discover the congestion at traffic level, which is based on the measure of congestion quantity at relay hop level node.

B. Measuring degree of blockage at Relay hop level node

Unlike conventional networks, nodes in the ad hoc network display a high level of heterogeneity with regard to both equipment and software configurations. Barrier capability and the heterogeneity of the exchange hop nodes can represent as diverse radio range, maximum retransmission counts. Hence the degree of stream loading, packet drop rate, and degree of buffer ingestion at relay jump degree node is minimum mix to obtain the level of congestion. The usage of these three purposeful beliefs supports to decouple the congestion measure process from other MLZBCC mechanism behavior.

The level of package drop rate, channel weight and amount of stream operation together provide a range to envisage the blocking due to improper ratio between retransmission and collision count. When retransmissions when compared with collision rate are significantly low then egress delay of relay hop node will increase proportionally, which leads to congestion and mirrored as congestion on account of buffer overflow.

C. Measuring degree of congestion at path level traffic

The degree of congestion at each relay hop together helps to identify the degree of congestion at path level traffic from supply to goal node. Each exchange hop amount node receives the degree of obstruction from its threshold designer. Since the destination node that will be last node of the course-plotting path isn't egress the status. Thus the destination node sounds to assess the amount of congestion at path stage traffic. Ergo to save the vitality, the congestion upgrade technique thinks two conditional actions, which follows:

1. Degree of blocking $d_c(h_i)$ at relay hop level node h_i will be send to its successor h_{i+1} if the '$d_c(h_i)$' is superior to the node level blockage threshold $d_c(\tau)$. Hence the energy conserve due to conditional transmission.
2. If degree of overcrowding at path level traffic $d_c(rp)$ that established by node h_i from its doorway initiator h_{i-1} is slighter than $d_c(h_i)$ then it update the $d_c(rp)$ else it remains same, hence energy protect due to prevention of $d_c(rp)$ update.

III. MULTIPLE MECHANISM CONGESTION CONTROL MODEL

The packet reducing often occurs in Manets. The reasons for this package reducing are as below

- Transmission Link failure.
- Inferred Transmission due to weigh down Ingress that leads Ingress getting power to low. This also can claim as packet reducing due to blocking at routing.

The congestion control can be evaluate in two stages by revolving over of the zonal head with the network partitioned into Cells as follows

- The Status of blocking at intra Cell level
- The status of jamming at inter Cell level

This helps in minimization of source level egress instruction cost and balances the power utilization.

TABLE1 NOTATIONS USED IN PROPOSED MODEL

Zone	**A geographical area, which is the part of preferred mobile ad hoc network**
ZCEA	Zone level congestion Evaluation Algorithm
ZERA	Zone level Egress Regularization Algorithm
ERA	Egress Regularization Algorithm
DPG	Distance Power Gradient
EIL	Ingress inferred Loss
LFL	Link Failure Loss
IRS	Ingress receiving strength
IRS_p	Present Ingress receiving strength

Zone	**A geographical area, which is the part of preferred mobile ad hoc network**
IRS_e	Expected Ingress Receiving Strength
RP	Routing Path
dt_n	Delay time at node n
N	Number of nodes in entire network
Zn_i	Number of nodes in a Cell i
zh_i	Cell head of the i^{th} Cell
zh'_i	Reserved Cell head of the i^{th} Cell
Z_c	Current Cell in the hierarchy
Z_p	Preceding Cell to the current Cell Z_c in hierarchy
Z_f	Fallowing Cell to the current Cell Z_c in hierarchy
Z_i	i^{th} Cell in the routing path
n_z	Cell of the node n
ζ_Z	Cell level Transmission Load Threshold
ζ_n	Node level Transmission Load Threshold
ζ_T	Predefined threshold that represents interval between two transmissions at one hop level
ζ_t	Actual interval between last two transmissions
ζ_{et}	Elapsed time since last transmission at one hop level
IRS_{ζ_T}	Average Ingress receiving strength threshold observed for predefined interval ζ_T
∂'	Average slopping threshold of the receiving strength
IRS_{ce}	Expected Ingress receiving strength threshold at current interval
IRS_r	Ingress receiving strength ratio
IRS_{cr}	Current ingress receiving strength ratio
BT_n	Buffering time at node n
$zdil_i$	Cell level degree of ingress load, here i is a Cell id.
$ndil_k$	Node level degree of ingress load, here k is the node id of Cell i

A. *Network and Node activities under projected protocol*

The network is to be crack into Cell For each Cell i where $i = 1..|Z|$; ($|Z|$ is total number of Cells)

Select Cell-head for each Cell i

Find spread load threshold ζ_n for each Cell i

By using ζ_n of each Cell extend load threshold for entire network can be measured.

B. *Splitting the network in to Cells*

We elect to the strategy described by Mohammad M. Qabajeh [8]. With all the understanding of the nodes the region is broken into equivalent partitions. Hexagon is mostly chased for the form for the reason that it covers a surface as well as offers the development of communication with more neighbors because they have near round shape of the transmitter. The availability of small, economical low power GPS receiver makes it possible to implement position-based in MANETs. The communication range of node is the side of hexagon along with denote as L. As the nodes ought to have the ability to correspond with one another the R and L are related as L=R/2.

Each Cell has a Cell characteristics (zid), Cell Header (zh) and Cell Leader Backup (zh'). The zh node maintains in succession about all the nodes in a Cell with their location and IDs. Also, maintain in sequence about the zh of the neighboring Cells as shown in the figure 1. The CLB node keeps a copy of the information stored at the zh so that it is not lost when the zh node is off or moving the Cell. By knowing the coordinates of a node position, nodes can achieve our self-mapping algorithm of their physical locations onto the current Cell and estimate it's zid easily. Figure 1.shows the general overview of the network architecture.

1) *Selecting Cell-Heads*

A Cell-Head selection happen under the pressure of the Subsequent metrics:

- Node positions: A node with a location *p* that is close to the centre is more likely to act as a Cell head.
- Optimum energy available: a node with higher energy e most likely acts as a Cell head.
- Computational ability: the node with high computational ability c is more potential to act as a Cell Head.
- Low mobility: the mobility *m of* a node is inversely comparative to its selection as a Cell head.

Each node of the Cell broadcasts its (p,e,c,m). The node that recognized itself as most optimal in (p,e,c,m) metrics, broadcast itself as Cell head zh. The next finest node in sequence claims itself as reserve Cell head zh'.

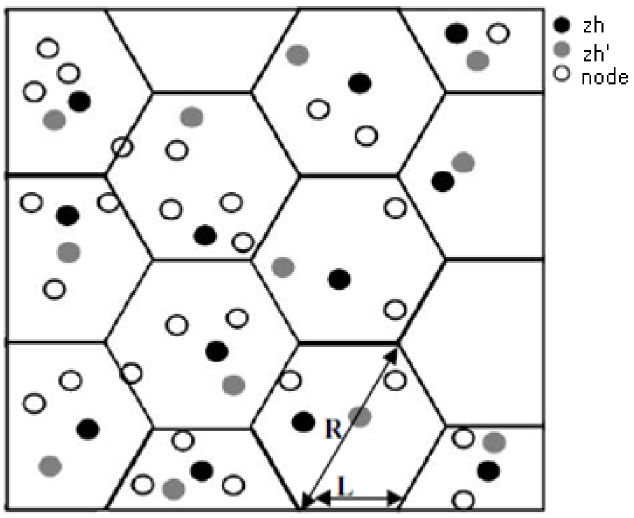

Fig 1: General overview of the Cell partitions in network

2) Information sharing at intra Cell level [Between Node and Cell head]

Each node n that is a subset to Cell Z verifies the Ingress load and distributes degree of ingress load dil_n with Cell head. Once $ndil_k$ received from each node k of the Cell i, the Cell head zh calculates the degree of ingress load at Cell level $zdil_i$.

$$zdil_{z_i} = \frac{\sum_{k=1}^{zn_i} ndil_k}{zn_i}$$

C. Zone level Congestion Evaluation Algorithm (ZCEA)

Zone level Congestion Evaluation Algorithm abbreviated as ZCEA is accessible in this section. ZCEA is an optimal algorithm that helps in locates the packet dropping under congestion. This assessment occurs under MLZBCC mechanism and then alerts network mechanism.

At an event of ingress receiving by node i :

Updating Ingress receiving strength:

$if\,(\zeta_t < \zeta_T)\,do$

$$\delta' := \frac{1}{2}\left(\frac{IRS_{cr} - IRS_{\zeta_T}}{\zeta_t}\right) + \frac{1}{2}(\delta')$$

$$IRS_{\zeta_T} := IRS_{cr}\left(\frac{\zeta_t}{\zeta_T}\right) + IRS_{\zeta_T}\left(\frac{\zeta_T - \zeta_t}{\zeta_T}\right)$$

$endif$

$if\,((\zeta_t) \neg < (\zeta_T))\,do$

$$\delta' := \frac{IRS_{cr} - IRS_{\zeta_T}}{\zeta_t}$$

$IRS_{\zeta_T} := IRS_{cr}$

$endif$

Detecting packet drop at MLZBCC mechanism level

$IRS_{ce} = IRS_{\zeta_T} + \delta' \zeta_{et}$

$if\,(IRS_{ce} < IRS_r)\,do$

$macAlert: link-failure$

$else$

$MacAlert: congestion$

$endif$

Fig2: ZCEA for determining congestion caused packet dropping

D. Zone level Egress Regularization Algorithm (ZERA)

This event occurs if MLZBCC-mechanism alert point out the congestion situation. Once the routing protocol [13] gets an alert from the MLZBCC mechanism a propos the blocking at a node i, it alerts the beneficiary citizen node which is the source node s for conflict node i. Hence s evaluates it's dil_s by evaluate with $zdil$ of Z_c (Cell of the node s). If dil_s is more in magnitude than $zdil_{z_c}$ the variation among dil_s and $zdil_{s_z}$ should be either greater or equal to the egress threshold ε then node s regularizes the egress load by influence its buffer time BT_s such that $ndil_s \geq zdil_{s_z} + \varepsilon_{s_z}$.

Here ε can be considered with following equation

$$\varepsilon_j = \frac{\sum_{k=1}^{zn_j} zdil_j - dil_k}{zn_j}$$

In case that the node s not able to normalize its egress so that disagreement node i terminates blocking then it alerts the zh_{s_z} (Cell-head of the Z_c , $s \in Z_c$). Subsequent that event zh_{z_c} alerts all the nodes in the network building the all nodes in the upstream of source node to way out load using the above stated slant. Then all nodes inform their $ndil$ and send to Cell-head zh_{z_c} , then Cell-head zh_{z_c} compute $zdil$ and confirms integrity of the $zdil$ by evaluation with dil . $zdil_{Z_c} \geq dil + \overline{\varepsilon}$ Concludes that congestion at contention node maintained by egress regularization at current Cell level. If $zdil_{z_c} < dil + \overline{\varepsilon}$ then CEA will be started at Z_p, which is adjacent upstream Cell to Z_c in contagious. In this process Cell head of the Z_c firstly alerts the Cell head of the counterpart Z_p then zh_{z_p} alerts all nodes that belongs to Z_p, of the route path. The above method of egress regularization at Cell level can be referred as ZERA (Zone level Egress Regularization Algorithm). Hence the nodes belong to Z_p normalize their egress load by utilize ZERA and alert Cell-head about their efficient degree of ingress load $ndil$. Then zh_{z_p} measures $zdil_{z_p}$ and verifies the consequence of $zdil_{Z_p} \geq dil + \overline{\varepsilon}$ .True indicates the elimination or minimization of congestion at the Cell due to the egress regularization at Cell Z_p, if false then Cell head of the Z_p performs the action of alert all other Cell heads using a broadcasting[12] apparatus about the congestion at adjacent Cell in downstream of the heridetary. Hence all Cells in the upstream side of the Z_p apply ZERA and the Cells in downstream side of the Z_p fill in there $zdil$. Then all Cells broadcast $zdil$ to resource Cell. Hence the source Cell revaluates the dil .Basing on the dil ,source node normalize its egress load.

Notations used in Algorithm:
i: Node that had been effected by emptiness
s: source node of the i.
Z_c : current Cell where $i, s \in Z_c$
Z_p : Immediate Cell to Z_c in upstream side of the pecking order.
$\{n_{u1}, n_{u2}, ..., n_{uk}\}_{Z_c}$: All upstream nodes to s .
$\{n_{d1}, n_{d2}, ..., n_{dk}\}_{Z_c}$: All downstream nodes to s .
$\{Z_S, Z_{u1}, Z_{u2}, ..., Z_{uk}\}$: Set of upstream Cells to Z_p in routing path, here Z_S is a Cell that contains source node of the routing path
$\{Z_{d1}, Z_{d2}, ..., Z_{dm}, ..., Z_T\}$: Set of downstream Cells to Z_p in routing path, here Z_T is a Cell that contain target node of the routing path
ε : Cell level egress threshold
$\overline{\varepsilon}$: Network level Egress threshold
Algorithm:
MLZBCC mechanism alerts about the blocking at node of Cell Z_c to routing protocol, hence the following steps perform in sequence

$$\varepsilon_{Z_c} = \frac{\sum_{k=1}^{zn_{Z_c}} zdil_{Z_c} - dil_k}{zn_{Z_c}}$$

complete following at node s
If $ndil_s > zdil_{Z_c}$ and $ndil_s - zdil_{Z_c} \geq \varepsilon_{Z_c}$ begin

$$BT_s = BT_s + bt$$

Note: Value of buffer threshold bt *should be certain such that* $dil_S \geq zdil_{Z_C} + \varepsilon_{Z_C}$
Return.
Endif
s Sends alert to zh_{Z_c} about conflict node i .
zh_{Z_c} alerts all nodes that belongs to Cell Z_c
$\{n_{u1}, n_{u2}, ..., n_{uk}\}_{Z_c}$ updates their $ndil$ by apply ZERA recursively and alerts zh_{Z_c}
$\{n_{d1}, n_{d2}, ..., n_{dk}\}_{Z_c}$ measures their $ndil$ and alerts zh_{Z_c}
zh_{Z_c} Measures $zdil$ as fallows

$$zdil_{Z_C} = \frac{\sum_{k=1}^{zn_{Z_C}} ndil_k}{zn_{Z_C}}$$

If $zdil_{Z_c} > dil$ and $(zdil_{Z_c} - dil) \geq \overline{\varepsilon}$ begin
Alert: blocking at contention node handle at current Cell Z_c level.
Return.
Endif
zh_{Z_c} Alerts zh_{Z_p}
zh_{Zp} Alerts all nodes that belong to Cell Z_p
For each node $n \in Z_p$ begin
If $ndil_n > zdil_{Z_p}$ and $ndil_n - zdil_{Z_p} \geq \varepsilon_{Z_p}$ begin

$$BT_n = BT_n + bt$$

Note: Value of barrier threshold bt *should be decided such that* $dil_n \geq zdil_{Z_C} + \varepsilon_{Z_C}$
Endif
Find dil_n *and send* dil_n *to* zh_{Z_p}
End-of-for each
zh_{Z_p} measures $zdil_{Z_p}$
if $zdil_{Z_p} > dil$ and $(zdil_{Z_p} - dil) \geq \overline{\varepsilon}$ begin
Alert: Egress regularization at Z_p leads to overcome congestion situation at contention Cell.
Return;
Endif
zh_{Z_p} Alerts all Cell heads in network regarding congestion contention Cell.

For each Cell z in $\{Z_S, Z_{u1}, Z_{u2}, ..., Z_{uk}\}$ begin
zh_z Alerts all nodes that belongs to Cell z
For each node $n \in z$ begin
If $ndil_n > zdil_z$ and $ndil_n - zdil_z \geq \varepsilon_z$ begin

$$BT_n = BT_n + bt$$

Note: Value of barrier threshold bt *should be understood such that* $dil_n \geq zdil_Z + \varepsilon_Z$
Endif
Find dil_n *and send* dil_n *to* zh_z
End-of-foreach
zh_z Measures $zdil_z$ and broadcast towards source Cell.
End-of-foreach

For each Cell z in $\{Z_{d1}, Z_{d2}, ..., Z_{dm}, ..., Z_T\}$ begin
For each node n belong to Cell z begin
determine $ndil_n$ and sends to zh_z
End-of-foreach
zh_z measures $zdil_z$ as

$$zdil_Z = \frac{\sum_{k=1}^{zn_Z} ndil_k}{zn_Z}$$

zh_z Sends $zdil_z$ to source Cell via propagation [12]
End-of-foreach
Z_S Measures dil as

$$dil = \frac{\sum_{i=1}^{|Z|} zdil_i}{|Z|}$$

Hence source node S of Cell Z_S, which is source node of the routing path standardize it's egress load to direction-finding path.

Fig 3: Zone level Egress Regularization Algorithm

IV. SIMULATIONS AND RESULTS DISCUSSION

In this section we discuss the outcome acquired from simulation conducted using “Madhoc Simulator” [16] in this section. We estimate concert using madhoc with the following considerations:

TABLE 2

Parameter used in Madhoc [16] for performance analysis

No of Hops:	225
Approximate Hop distance	300 meters
Approximate total network	1000X1000 meters
fairly accurate zone Radius	100X100 meters
Physical channel bandwidth	2mbps
MLZBCC Mechanism	802.11 DCF with option of handshaking prior to data transferring
Physical mechanism illustration	802:11B
presentation Index	Egress regularization cost and end-to-end throughput
Be very successful simulation time	150 sec

The replications are conducted on three routes opposed by the no of hops and length.

1. Short length path: A route with 15 hops
2. middling length : A route with 40 hops
3. Max Length: A route with 81 hops

Exactly the same load is given to all the paths having a regular interval of 10 sec. The fig 5 indicates the improvement of CDC-CPF over congestion control protocol [15] in impediment control cost. A. The congestion detection cost examination between CDC-CPF and congestion control protocol [15] is explored in the energy good organization that is elevated by fig 6 realized under CDC-CPF.

The method of capacity of congestion get a handle on and congestion diagnosis charge is as follows:

Based on the reference ease of use, bandwidth and power, for specific operation a threshold value between 0 and 1 assigned. In the system of congestion detection and get a grip on the total cost is determined by summing the cost limit of every event. In fig 5 the judgment between obstructions expenses seen for CDC-CPF and blocking and rivalry get a handle on type [15] are shown.

$$\cos t_{ch} = \sum_{e=1}^{E} ct_e$$

Here $\cos t_{ch}$ is the price of a blockage controlling activity ch , E is total amount of events included. ct_e Is the threshold cost of an event e .

The example events are:

1.” alert to source node from MLZBCC mechanism”
2. “Alert from node to Cell head”, “propagation by Cell head to other Cell heads”
3. “Ingress judgment and egress regularization”.
4. Alert about $d_c(h_i)$
5. Bring up to date $d_c(rp)$

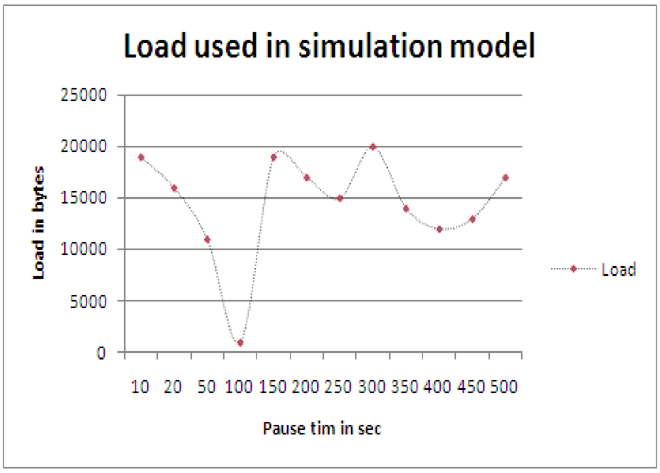

Fig 4: Load in bytes drive by source node of the routing path [In regular interval of 10 sec]

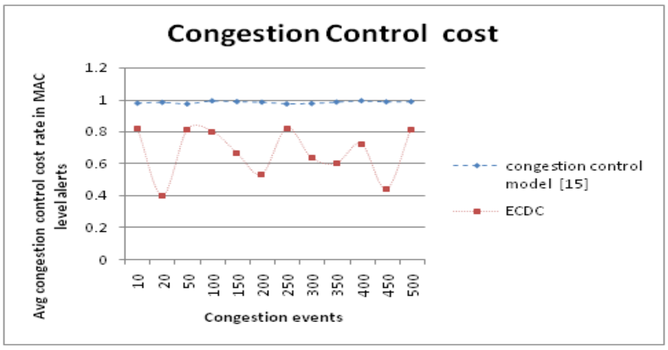

Fig 5: Congestion Control cost comparison chart

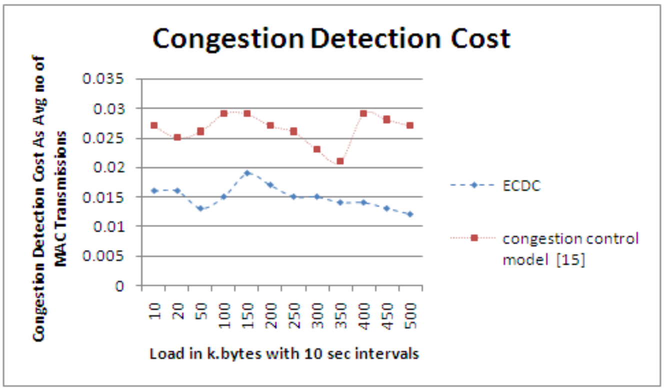

Fig 6: A line chart comparison of Congestion detection cost

V. Conclusion

This manuscript discussed about suggested "A Multi Level Zone Based Congestion Control Mechanism for Multicast Routing in MANETS" suggests source level Egress regularization cost can be minimized and energy utilization for congestion status alerts can also be balanced

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Authors:

A.Mallaiah received the M.E degree in Electronic Instrumentation from Andhra University, Visakhapatnam in 2004, B.Tech degree in Electronics and communication Engineering from R.V.R & J.C College of Engineering, Guntur in 2002. He is Associate Professor, Department of Electronics and Communication Engineering at Gudlavalleru Engineering College, Gudlavalleru. He has a total teaching experience (UG and PG) of 10 years. He has guided and co-guided 6 P.G students. He is presently doing Ph.D from JNTU Ananthapur. His research areas include Embedded Systems, Molecular Electronics, Digital Signal Processing and Digital Image Processing.

P.Imran Khan is pursuing his M.Tech degree in Digital Electronics & Communication Systems from Gudlavalleru Engineering College, Gudlavalleru, B.Tech from Electronics & Communication Engineering from Nimra College of Engineering and Technology, Vijayawada in 2011. His research interest includes mobile adhoc networks and Embedded Systems.

Dr.G.N.Swamy obtained his Bachelor’s degree from Acharya Nagarjuna University, Master’s degree from BHU &Ph.D from Andhra University. Areas of interest are Digital signal processing, Digital Image Processing, Electronic Circuits and Microwave Engineering He a having experience of 22 years in the field of teaching. He published 35techinical papers in national and international journals and conferences. He reviewed number of papers for international journal and international conferences. He is a Fellow of IETE and IE. Presently working as Professor& HOD of E&I Dept., in V.R Siddhartha Engineering College, Vijayawada, India.